

# Intel Labs Haskell Research Compiler

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- does whole program compilation;
- achieves overall performance parity with GHC+LLVM;
- is significantly better for some programs;
- does automatic SIMD vectorization for Intel CPUs.

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Known Limitations:

- No lightweight threads, sparks, or STM (easy);
- No exception re-throw for thunks (fixable);
- No asynchronous exceptions (hard).

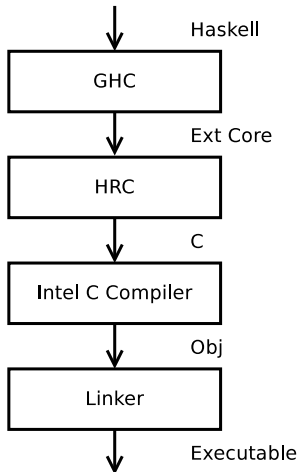
# Functionality

HRC is highly compatible to GHC:

- Modified GHC and base libraries to handle differences;
- Modified Vector library to use initializing writes;
- Modified Cabal to compile for HRC;
- Little or no modifications to most other libraries.

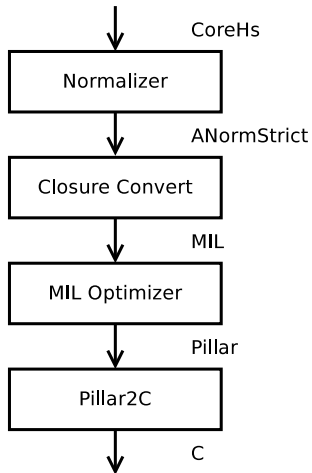


# Compilation Process



# Inside HRC Pipeline

CoreHs	functional, non-strict
ANormStrict	ANF, strict explicit thunk/eval
MIL	CFG/SSA based high-level object
Pillar	inspired by C-- variant of C



# Comparison to GHC

GHC		HRC	
Desugaring Type analysis Core-to-Core transformation			
STG	Functional, object memory model, optimized for currying and thunks	MIL	CFG/SSA based, object memory model, conventional
Cmm	CFG blocks, low-level types, and custom calling convention	Pillar	C types, C calling convention meta and GC support
LLVM or NCG	Portable LLVM bitcode, or direct assembly generation	Intel C/C++ Compiler	Portable C code compiled to assembly
Runtime and GC	Optimized for currying and thunks	Runtime and GC	Conventional

# MIL

- High level object model with low level control flow
- Leveraging immutability and memory safety
- Data flow and control flow analysis
- Inter- and intra- procedural optimizations
- Representation/contification/loop/thunk optimizations
- SIMD auto vectorization

# Immutable Array

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- creating a mutable array;
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- create immutable array
- do initializing writes
- read

Programmers must ensure:

- an array field is written to before it is read;
- a field is never written to more than once.

# Walkthrough by Example

`odd, even :: Int → Bool`

`odd 0 = False`

`odd n = even (n - 1)`

`even 0 = True`

`even n = odd (n - 1)`



# GHC Core

```
even :: (Int → Bool) =
  \ (n :: Int) →
    %case n %of (- :: Int)
      {I# (n1 :: Int#) →
        %case n1 %of (n2 :: Int#)
          {(0 :: Int#) → True;
           %_ → odd (I# (n2 - (1 :: Int#))))}};
```

```
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# GHC Core

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```

## ANormStrict

```
even = %thunk
  %let f = \n →
    %let n0 = %eval n
    %in %case n0 of
      {I# n1 → %case n1 of
        {0 → %eval True;
          - → %let u = %let v = %let w = 1
                                %in n1 - w
                                i = %eval I#
                                %in i v
                                t = %thunk u
                                g = %eval odd
                                %in g t}}
      %in f;
```

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even = %thunk
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                    %in i v
                    t = %thunk u
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                    %in g t}}
```

```
%in f;
```



## ANormStrict (closure converted)

```
even = %thunk <I#, True, odd;>
  %let f = \<I#, True, odd; n> →
    %let n0 = %eval n
    %in %case n0 of
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        {0 → %eval True;
         _ → %let u = %let v = %let w = 1
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```

# MIL

```
f_code =  
  Code(CcClosure(lv_I#, lv_True, lv_odd); n)  
  {  
    L0():  n0 ← Eval(n) ⇒ L1  
    L1():  Case tagof(n0) { U32(0) ⇒ L2() }  
    L2():  n1 = SumProj(n0.U32(0).0)  
           Case n1 { S32(0) ⇒ L8() Default ⇒ L3() }  
    L3():  v = SInt32Minus(n1, 1)  
           i ← Eval(lv_I#) ⇒ L4  
    L4():  u ← CallClos(i) (v) ⇒ L5  
    L5():  t = ThunkMkVal(u)  
           g ← Eval(lv_odd) ⇒ L6  
    L6():  c ← CallClos(g) (t) ⇒ L7  
    L7():  Goto L10(c)  
    L8():  b ← Eval(lv_True) ⇒ L9  
    L9():  Goto L10(b)  
    L10(r): Return(r)  
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          e = SInt32Eq (n1, 0)  
          Case e { True ⇒ L6() False ⇒ L2() }  
    L2():  v = SInt32Minus(n1, 1)  
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  }
```

## MIL (after Rep optimization)

```
f_code =  
  Code(CcCode; n)  
  {  
    L0():  even = SInt32Eq (n, 0)  
           Case even { True ⇒ L4() False ⇒ L1() }  
    L1():  v = SInt32Minus(n, 1)  
           odd = SInt32Eq(v, 0)  
           Case odd { True ⇒ L3() False ⇒ L2() }  
    L2():  w = SInt32Minus(v, 1)  
           Call(f_code) (w) =  
    L3():  Return(gv_False)  
    L4():  Return(gv_True)  
  }
```

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f_code =  
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           odd = SInt32Eq(v, 0)  
           Case odd { True ⇒ L3() False ⇒ L2() }  
    L2():  w = SInt32Minus(v, 1)  
           Call(f_code) (w) ⇒  
    L3():  Return(gv_False)  
    L4():  Return(gv_True)  
  }
```



## MIL (after contification)

```
f_code =  
Code(CcCode; n)  
{  
  L0(): Goto L1(n)  
  L1(u): even = SInt32Eq(u, 0)  
         Case even { True ⇒ L5() False ⇒ L2() }  
  L2(): v = SInt32Minus(u, 1)  
         odd = SInt32Eq(v, 0)  
         Case odd { True ⇒ L4() False ⇒ L3() }  
  L3(): w = SInt32Minus(v, 1)  
         Goto L1(w)  
  L4(): Return(gv.False)  
  L5(): Return(gv.True)  
}
```

## MIL (after arithmetic simplification)

```
f_code =  
  Code(CcCode; n)  
  {  
    L0():  Goto L1(n)  
    L1(u):  even = SInt32Eq(u, 0)  
           Case even { True ⇒ L5() False ⇒ L2() }  
    L2():  odd = SInt32Eq(u, 1)  
          Case odd { True ⇒ L4() False ⇒ L3() }  
    L3():  w = SInt32Minus(u, 2)  
          Goto L1(w)  
    L4():  Return(gv_False)  
    L5():  Return(gv_True)  
  }
```

# Benchmarking

40+ benchmark programs:

- a mixed set of programs from nofib benchmark suite;
- performance oriented programs using array libraries.

Sequential performance is measured by compiling and running on 2.7GHz Xeon machine (32-bit Windows) with:

- standard GHC 7.6.1
- GHC 7.6.1 + LLVM 2.9
- HRC with modified GHC 7.6.1 + Intel C/C++ compiler.

# Benchmark Result

HRC is at parity to GHC+LLVM, which is 10% faster than GHC.

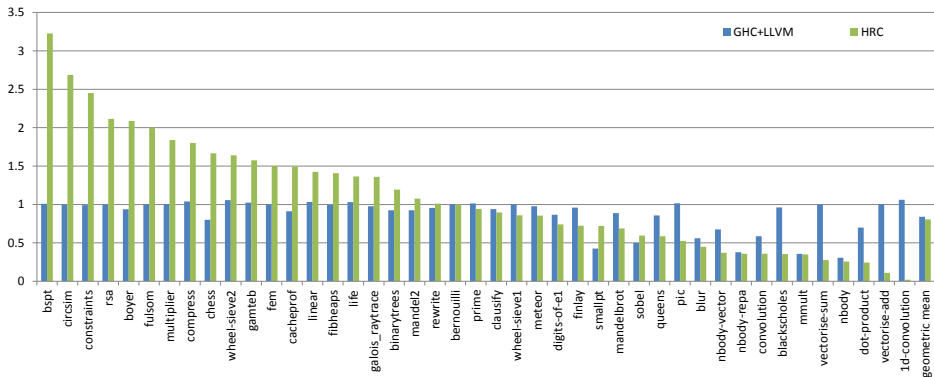


Figure: Kernel Execution Time Relative to GHC (smaller is better)

## Benchmark Result (selected)

Performance oriented program with a numeric computation kernel using arrays.

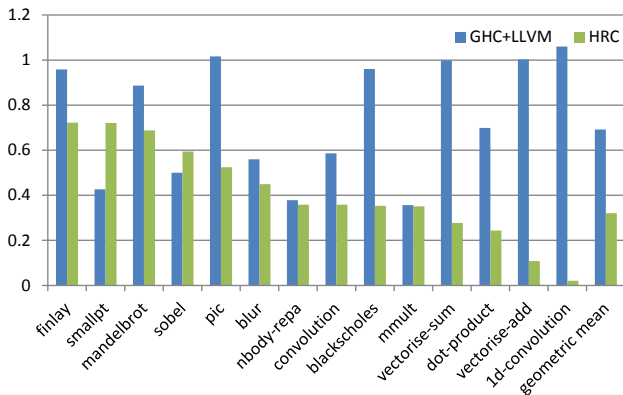


Figure: Kernel Execution Time Relative to GHC (smaller is better)

# Performance

For GHC:

- GHC is better at executing lazy code and curried functions.
- GHC's object representation and GC are better suited to typical Haskell programs.

For HRC:

- HRC focuses on optimizing strict programs with hot loops.
- HRC benefits significantly from the elimination of thunks, boxes and branches.

# Take Aways

- Reusing GHC is a big win (Core: easy, library: a bit work)
- Novel MIL design choices:
  - Low-level control with high-level object.
  - Immutable array with initializing writes.
- Eliminating thunks is critical to performance.
- Penalty of not using a specialized runtime.
- Compilation through C has overhead, but not as significant.

## References

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